Semantics 1

May 8, 2012

Gerhard Jäger



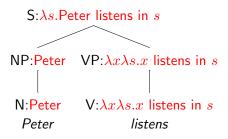


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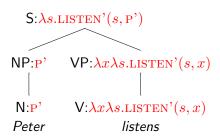
- sentence meaning = lexical meaning + syntax
- example:

Peter listens

- sentence meaning: λs .Peter listens in s
- lexical meanings:
 - ||Peter|| = Peter
 - $\| \textit{listens} \| = \lambda x \lambda s. x \text{ listens in } s$
- syntax: [S[NP[N] Peter]][VP[V] Iistens]]]



- ullet So far, we used English + some lambda notation as meta language.
- Predicate logic is more precise than English; therefore it is to be preferred as meta language.
- note: all predicates have an additional argument for situations. (This is different from the translations you used in your logics class.)



- meaning of the mother node can be computed from the meanings of the daughter nodes:
 - for non-branching nodes, mother node and daughter node have the same meaning
 - in an NP-VP structure, the meaning of the VP (which is a function) is applied to the meaning of the NP
- Assumption: this correspondence between syntax and semantics holds for all English sentences. (The correct syntax of English is of course much more complex, but I try to keep things simple for expository purposes.)

- formally: for each syntactic rule, there is a corresponding semantic rule
- so far, we have
 - $S \to NP, VP :: ||S|| = ||VP||(||NP||)$
 - $NP \to N :: ||NP|| = ||N||$
 - $\bullet \ VP \to V :: \|VP\| = \|V\|$

Schönfinkeling (a.k.a. Currying)

- meaning of transitive verb: two-place relation
- e.g.: *loves* $\sim \{\langle x,y \rangle | \text{LOVE'}(x,y) \}^1$
- expression as characteristic function:

$$\lambda \langle x, y \rangle \in E \times E.$$
LOVE' (x, y)

lambda conversion:

$$(\lambda \langle x, y \rangle \in E \times E.$$
LOVE' $(x, y))(\langle a, h \rangle) = LOVE'(a, h)$

¹We ignore situation dependence for a moment.

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Schönfinkeling

• What is the meaning of *loves John*? The set of individuals that love John.

$$\|loves\ John\| = \{x| \text{LOVE'}(x,j)\} \approx \lambda x. \text{LOVE'}(x,j)$$

• *loves* can also be considered as a function that maps the meaning of α to the meaning of *loves* α :

$$\|loves\| = \lambda y \lambda x. LOVE'(x, y)$$

Schönfinkeling

• two-place relation $\{\langle x,y\rangle | \text{LOVE'}(x,y) \}$ is transformed into two-place characteristic function $\lambda \langle x,y\rangle. \text{LOVE'}(x,y)$, which, in turn, can be transformed into a one-place function with a one-place characteristic function as its value:

$$\lambda y \lambda x.$$
LOVE' (x, y)

• general recipe:

$$\{\langle x,y\rangle|R(x,y)\} \rightsquigarrow \lambda\langle x,y\rangle.R(x,y) \rightsquigarrow \lambda y\lambda x.R(x,y)$$

• same principle also applies to *n*-ary relations:

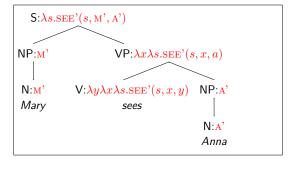
$$\{\langle x_1, \cdots, x_n \rangle | S(x_1, \cdots, x_n)\} \rightsquigarrow \lambda x_n \cdots \lambda x_1 \cdot S(x_1, \cdots, x_n)$$

Note: Order of the variables in the λ -prefix is mirror image of their order within the argument frame of the relation!

Transitive Verbs

- examples: love, know, see, help, ...
- express two-place relations between individuals
- if situation dependence is added, we get three-place relations
- $\| \textit{Mary sees Anna} \| = \lambda s.\text{SEE'}(s, \text{M'}, \text{A'})$
- $\|sees\| = \lambda y \lambda x \lambda s. SEE'(s, x, y)$

Transitive Verbs



Rules:

- $\bullet \ S \to NP, VP :: \\ \|S\| = \|VP\|(\|NP\|)$
- $\bullet \ NP \rightarrow N :: \\ \|NP\| = \|N\|$
- $VP \to V ::$ ||VP|| = ||V||
- $\bullet \ VP \rightarrow V, NP :: \\ \|VP\| = \|V\| (\|NP\|)$

The compositional analysis of the Boolean operators can also be expressed in this format:

Negation

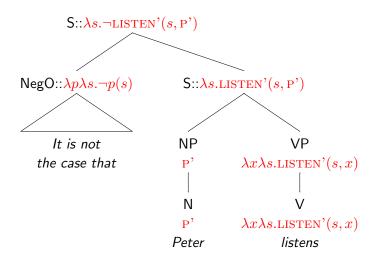
- Logical operator of negation can be expressed in two ways in English:
 - It is not the case that Peter listens.
 - Peter doesn't listen.
- in both cases, the semantic effect is set complementation:

 $\|Peter\ does\ not\ listen\| = \lambda s. \neg \texttt{LISTEN'}(s,p)$

Negation

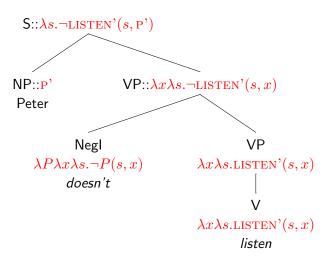
- New rules:
 - $S_1 \to NegO, S_2 :: ||S_1|| = ||NegO||(||S_2||)||$
 - $VP_1 \to NegI, VP_2 :: ||VP_1|| = ||NegI|| (||VP_2||)||$
 - $NegO \rightarrow It$ is not the case that :: $\|NegO\| = \lambda p\lambda s. \neg p(s)$
 - $NegI \rightarrow doesn't :: \|NegI\| = \lambda P \lambda x \lambda s. \neg P(x,s)$

Negation



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Negation

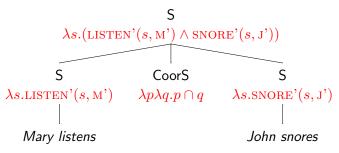


Sentence Coordination

- Rules:
 - $S_1 \to S_2, CoorS, S_3 :: ||S_1|| = ||CoorS||(||S_2||)(||S_3||)$
 - $CoorS \rightarrow and :: \lambda p \lambda q. p \cap q$
 - $CoorS \rightarrow or :: \lambda p \lambda q. p \cup q$
- Note:

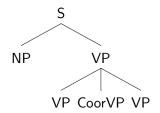
$$\lambda s.\phi \cap \lambda s.\psi = \lambda s.(\phi \wedge \psi)$$
$$\lambda s.\phi \cup \lambda s.\psi = \lambda s.(\phi \vee \psi)$$

Sentence coordination



VP coordination

- Coordination may conjoin two VPs
 - Peter sleeps and snores.
 - John walks and talks.
- syntactic structure:



semantics: similar to sentence operators
Peter sleeps and snores ⇔ Peter sleeps and Peter snores.

VP coordination

- Rules:
 - $VP_1 \to VP_2, CoorVP, VP_3 :: ||VP_1|| = ||CoorVP||(||VP_2||)(||VP_3||)$
 - $CoorVP \rightarrow and :: \lambda P \lambda Q \lambda x \lambda s. P(x)(s) \wedge Q(x)(s)$
 - $CoorVP \rightarrow or :: \lambda P \lambda Q \lambda x \lambda s. P(x)(s) \vee Q(x)(s)$

Boolsche Operatoren

VP coordination

