Semantics 1

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Pronouns

- Reflexive and personal pronouns depend on context for their interpretation.
- Reflexives can be interpreted in two ways:
- **O** coreferent:
- (1) Peter_i shaved himself_i.
 - (semantically) bound:
- (2) a. by a quantifier: Every player_i shaved himself_i
 - **b**. **by a wh-phrase:** A man who_i shaved himself_i arrived.

Pronouns

- Personal pronouns can be interpreted in three ways:
- **o** free: (denotation is fixed by the context of utterance)
- (1) Mary_i likes him_i.

Oreferent:

(3)

(2) [Peter_i's father]_k shaved him_i.

2 (semantically) bound:

- a. by a quantifier: Every $philosopher_i$ praised a book that he_i wrote.
- b. by a wh-phrase: A philosopher who_i praised every book that he_i wrote arrived.

• lexical meaning:

 $\|\mathsf{himself}_i\| = x_i$

(and likewise for all other reflexives)

- Coreferent reading:
- (1) Peter₁ shaved himself₁.
 - Index 1 on name *Peter* restricts the context of interpretation to such assignment functions where

$$x_1 = \mathbf{P}'$$

• compositional derivation:



• contextual equivalence:

 $x_1 = \mathbf{P'} \vdash \lambda s. \text{SHAVE'}(s, \mathbf{P'}, x_1) = \lambda s. \text{SHAVE'}(s, \mathbf{P'}, \mathbf{P'})$

- binding by a quantifier
- (1) Every player₁ shaved himself₁.



Recall the interpretation rule for the root node, ie. a structure that results from QR:

- $||S^{\uparrow}|| = ||NP_1||(\lambda x_1.||S^{\downarrow}||)$
 - $= \lambda Q \lambda s. \forall x (\text{player}'(s, x) \to Q(s, x)) (\lambda x_1 \lambda s. \text{shave}'(s, x_1, x_1))$
 - $= \lambda s. \forall x (\text{player'}(s, x) \rightarrow \text{shave'}(s, x, x))$

- Rule for QR involves applying λx_1 to the lower S-segment
- this λ-operator binds the (variable corresponding to) the trace as well as the (variable corresponding to) the reflexive
- Unlike in cases of coreference, the final interpretation does not contain free variables, and the indexing does not impose constraints on the context

- binding by a *wh*-pronoun
- (1) A man who shaved himself arrived.



Recall the interpretation rule for the \overline{S} -node, ie. a structure that results from *wh*-movement:

$$\begin{split} \|\overline{S}\| &= \|_1 \| (\lambda x_1 \| S \|) \\ &= \lambda P \lambda Q \lambda x \lambda s. Q(s, x) \wedge P(s, x) (\lambda x_1 \lambda s. \text{SHAVE}'(s, x_1, x_1)) \\ &= \lambda Q \lambda x \lambda s. Q(s, x) \wedge \text{SHAVE}'(s, x, x) \end{split}$$

- Rule for interpreting moved wh-elements involves applying λx_1 to the S-node
- this λ-operator binds the (variable corresponding to) the trace as well as the (variable corresponding to) the reflexive

• Free pronouns:

- coindexed neither with a c-commanding binder (quantifier or wh-phrase) nor with any proper noun within the same sentence
- correspond to free variables in semantic representation
- interpretation is determined by assignment function, i.e. by the context
- behave like proper nouns with respect to semantic composition

(1) John₁ shaved him₂



• Coreferent pronouns:

- coindexed with some proper noun within the same sentence
- due to Binding Principle B, this NP must not c-command the pronoun if it occurs within the same local clause
- behave like free pronouns with respect to semantic composition
- interpretation is constrained by context (just like for coreferent reflexives)

(1) [Every student from Tübingen₁]₂ likes it₁.



• Index 1 on name *Tübingen* restricts the context of interpretation to such assignment functions where

$$x_1 = T'$$

• contextual equivalence:

 $x_1 = \mathbf{T}' \vdash ||S|| = \lambda s. \forall x (\text{STUDENT}'(s, x) \land \text{FROM}'(s, x, \mathbf{T}') \to \text{Like}'(s, x, \mathbf{T}'))$

Bound pronouns:

- co-indexed with a binder i.e. a quantifier or a *wh*-phrase that **c-commands** the pronoun **at LF**
- due to Binding Principle B, this binder must **not c-command** the pronoun at **S-structure**

- (1) [A student from [every city]₁]₂ likes it₁.
 - S-Structure: no c-command ~> BT Principle B is fulfilled



• LF: inverse linking reading ~> c-command ~> binding



• LF: narrow scope reading \rightsquigarrow no c-command \rightsquigarrow pronoun remains free



Derivation of previous reading, step by step

$$\begin{split} \|PP^{\downarrow}\| &= \lambda P \lambda x \lambda s. P(s, x) \wedge \text{FROM}'(s, x, x_1) \\ \|PP^{\uparrow}\| &= \lambda P \lambda x \lambda s. P(s, x) \wedge \|NP_1\|(s, \lambda x_1 \lambda s. \|PP^{\downarrow}\|(\lambda x \lambda s. \top)(x)) \\ PP^{\downarrow}\|(\lambda x \lambda s. \top)(x) &= (\lambda P \lambda x \lambda s. P(s, x) \wedge \text{FROM}'(s, x, x_1))(\lambda x \lambda s. \top)(x) \\ &= \top \wedge \text{FROM}'(s, x, x_1) \\ &= \text{FROM}'(s, x, x_1) \\ \|PP^{\uparrow}\| &= \lambda P \lambda x \lambda s. P(s, x) \wedge \|NP_1\|(s, \lambda x_1 \lambda s. \text{FROM}'(s, x, x_1)) \\ &= \lambda P \lambda x \lambda s. P(s, x) \wedge \\ &\quad (\lambda Q \lambda s. \forall y(\text{CITY}'(s, y) \to Q(s, y)))(\lambda x_1 \lambda s. \text{FROM}'(s, x, x_1)(s) \\ &= \lambda P \lambda x \lambda s. P(s, x) \wedge \forall y(\text{CITY}'(s, y) \to \text{FROM}'(s, x, y))) \end{split}$$